Windows Into the Past: Exploiting Legacy Crypto in Modern OS Kerberos Implementation

> Michal Shagam Ronen Tel Aviv University



Eyal

# Our Findings

- We broke Kerberos!
- Improved the Bleichenbacher attack
- End-to-end attack on Windows
- Windows used everywhere



### What is the Kerberos Protocol?

- Authentication protocol over non-secure networks
- Provides Secure communication
- Provides Secure access to services



# Simplified Look At Kerberos



### In Passwords We Trust?

- Dictionary Attacks
- Phishing
- Password DB Breaches



## Simplified Look At Kerberos



### No More Passwords – Secure?

• We got rid of the password – Security!





# Bleichenbacher Attack Returns



- Bleichenbacher Attack (1998) PKCS#1 V1.5 vulnerable
- Manger's Attack (2001) both PKCS#1 V1.5 and V2.0 vulnerable
- More and more attacks (SSLv2, TLS):



• No attacks on Kerberos – Why?

# Possible Bleichenbacher Oracles in Kerberos

- Finding Bleichenbacher oracles in Kerberos is not trivial
- Error messages
  - no error messages sent connection times out
- Timing analysis
  - connection timeout independent of RSA decryption
- Microarchitectural side channels
  - RSA decryption is done on the smartcard



# Layers of RSA Decryption

- Modeled in layers
  - Mathematical computation
  - Unpadding Removal and validation of padding
  - Protocol level error handling and padding oracle mitigations
- Mathematical and data conversion runs on smartcard
- Unpadding and protocol level handling runs on Client CPU
  - This can be targeted by a microarchitectural side channel attack!





### Survey of Common Kerberos Implementations

- Windows and Heimdal (used by MacOS) nonconstant time protocol implementations, no Bleichenbacher mitigations
- MIT implementation implements mitigations
- Windows and OpenSC smartcard interfaces are nonconstant time

OS	OS Kerberos		Vulnerable
	Protocol	Interface	
Windows	Security Pack-	basecsp	both
	age		
macOS	Heimdal	OpenSC	both
FreeBSD	Heimdal	OpenSC	both
Ubuntu	MIT	$\operatorname{OpenSC}$	interface
RedHat	MIT	$\operatorname{OpenSC}$	interface
Gentoo	MIT	$\operatorname{OpenSC}$	interface
OpenSuse	MIT	$\operatorname{OpenSC}$	interface
CentOS	MIT	$\operatorname{OpenSC}$	interface
ArchLinux	MIT	$\operatorname{OpenSC}$	interface
Suse	MIT	OpenCryptoki	no <sup>a</sup>

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Implementation Vulnerability Survey

### **REing Windows**

### . F K W B - 🎕 🍇 🗠 🗠 🗸 🛱 🖄 📾 🗣 🏤 🛇 💷 🔶 🗐 😓 🧶

: Decompile: CspUnpadData - (basecsp.dll) 😵 🐂 📓 👻 🗙	Functions - 46 items (a	of 478 )		III 🖸 🚺	×	🖼 Listi	ing: basecsp.dll		Q.	d 💼   🖳 🖳 🎀 🎉 🐞   📑 • 🗙
· · · · · · · · · · · · · · · · · · ·	Label	Location	Function Signature	Function Size		base	ecsp.dll 🗙			
/* unsigned longcdecl CspUnpadData(struct _CARD_RSA_DECRYPT_INFO *ptr64,unsigned long *	InitializeCspState	180001c4c	ulong Initialize		288					1000063E4 (K)
ptr64,unsigned char *ptr64 *ptr64) */	InitializeCspCaching	18000ec9c	ulonglong Initia		177		undefined4	Stack[-0x38	]:4 local_38	XREF[1]: 180006475(W)
	I_CspReadCardCacheFile	18000e3b8	ulong I_CspReadC		316			?CspUnpadData@	@YAKPEAU_CARD_RSA_DECRYPT_INFO@@ XREF[5]	: FindCardBySerialNumber:]
ulong CspUnpadData(_CARD_RSA_DECRYPT_INFO *param_1,ulong *param_2,uchar **param_3)	I_CspQueryKeySizes	18000bb98	ulong I_CspQuery		604			CspUnpadData		FindCardBySerialNumber:]
	I_CspQueryFreeSpace	18000b86c	ulong I_CspQuery		592					LocalAcquireContext:1800
{	I_CspQueryCapabilities	18000a164	ulong I_CspQuery		587					LocalAcquireContext:1800
undefined uVarl;	I_CspCacheInitializeQuer.	18000e964	void I_CspCacheI		81					180023ab3(*)
ulong uVar2;	GetCspState	180001d74	ulong GetCspStat		82	@	180006330 48 89 5c	MOV	<pre>qword ptr [RSP + local_res10],RBX</pre>	
undefined8 uVar3;	DeleteCspState	180001908	ulong DeleteCspS		621		24 10			
char **ppcVar4;	CapWriteFile	180000970	undefined CspWri		437		180006335 55	PUSH	RBP	
char **ppcVar5;	CspWinter ine CspVerifiCachedPinConsi	18000505	undefined CopWar		275		180006336 56	PUSH	RSI	_
uint uVar6;	CoplingadData	180006330	wlong CepUppadDa		353		180006337 57	PUSH	RDI	
char **ppcVar7;	CspRemoveCachedPin	18000ab44	undefined CspRem		264		180006338 41 56	PUSH	R14	
uint uVar8;	CspReAllocH	1800 1a6e0	undefined CspReA		59		18000633a 41 57	PUSH	R15	_
ulong local_res8 [2];	CspReadUserCertificate	18000dea0	undefined CspRea		334		18000633c 48 83 ec 30	SUB	RSP, 0x30	
char **local_res20;	CspReadFile	18000b3ac	undefined CspRea		546		180006340 33 ff	XOR	EDI,EDI	
	CspQueryKeySizes	18000bdfc	ulonglong CspQue		214		180006342 4c 8b fa	MOV	R15,param_2	
ppcVar7 = (char **) 0x0;	CspQueryFreeSpace	18000bac4	ulonglong CspQue		205		180006345 33 d2	XOR	param_2,param_2	_
local_res8[0] = 0;	CspQueryCardCacheForIt.	18000e724	ulonglong CspQue		567		180006347 89 7c 24 60	MOV	dword ptr [RSP + local_res8],EDI	-
local_res20 = (char **) 0x0;	CspQueryCapabilities	18000a3b8	ulonglong CspQue		213		18000634b 8b f7	MOV	ESI,EDI	<b>E</b>
<pre>WppTraceIndent((char **)param_1,0);</pre>	CspPadData	18000bee0	ulonglong CspPad		147		18000634d 48 89 7c	MOV	<pre>gword ptr [RSP + local_res20],RDI</pre>	-
<pre>ppcVar4 = (char **)WPP_GLOBAL_Control;</pre>	CspLeaveCriticalSection	1800 1a7b8	undefined CspLea		87		24 78			2
if ((((undefined **)WPP_GLOBAL_Control != sWPP_GLOBAL_Control) ss	CspInitializeCriticalSection	n 18001a728	undefined8 CspIn		26		180006352 4d 8b f0	MOV	R14,param_3	
(ppcVar4 = (char **)WPP_GLOBAL_Control, (WPP_GLOBAL_Control[0x1c] & 2) != 0)) &&	CspIncrementCacheFresh.	18000e4fc	ulonglong CspInc		263		180006355 48 8b d9	MOV	RBX,param_1	
<pre>(ppcVar4 = (char **)WPP_GLOBAL_Control, 4 &lt; (byte)WPP_GLOBAL_Control[0x19])) {</pre>	CspImportSimpleBlobHelp	er 180005ba4	ulong CspImportS		292		180006358 e8 db ae	CALL	WppTraceIndent	void WppTraceIndent(
<pre>ppcVar4 = *(char ***) (WPP_GLOBAL_Control + 0x10);</pre>	CspGetContainerInfo	180008684	undefined CspGet		1126		ff ff			
WPP_SF_s(ppcVar4,0x3a,&WPP_e8978bfc35ec3505198246492c009beb_Traceguids,WPP_pszIndent);	CopEnterCriticalSection	180018750	undefined8 CspEn		169		18000635d 48 8b 0d	MOV	param_1, gword ptr [->WPP_GLOBAL_Control]	= 18002c0f8
3	CopDeleteFile	18000c05c	undefined Capbal		275		94 5d 02 00			
<pre>if (((param_1 == (_CARD_RSA_DECRYPT_INFO *)0x0)    (param_2 == (ulong *)0x0))   </pre>	CspDeleteContainer	18000a494	undefined Capbel		264		180006364 48 8d 2d	LEA	RBP, [WPP_GLOBAL_Control]	= 18002c0f8
(param_3 == (uchar **)0x0)) {	CspCreateFile	18000b2d4	ulonglong CspCre		208		8d 5d 02 00			
uVar2 = 0x57;	CspCreateContainer	18000a5a4	ulonglong CspCre		214		18000636b 48 3b cd	CMP	param_1,RBP	_
ppcVar5 = ppcVar7;	CspCheckContainerMapRe	18000c734	ulonglong CspChe		190		18000636e 74 26	JZ	LAB_180006396	
3	CspCacheLookupFileProc	18000eaf0	ulong CspCacheLo		287		180006370 f6 41 1c 02	TEST	byte ptr [param_1 + 0x1c]=>DAT_18002c114	,0x2
else {	CspCacheDeleteFileProc	18000ec20	ulong CspCacheDe		114		180006374 74 20	JZ	LAB_180006396	
if (*(int *)param_1 != 2) {	CspCacheAddFileProc	18000e9c0	ulong CspCacheAd		286		180006376 80 79 19 05	CMP	byte ptr [param_1 + 0x19]=>DAT_18002c111	,0x5 💻
uVar2 = 0x51a;	CspBuildUserCertStore	18000d500	ulonglong CspBui		1267		18000637a 72 la	JC	LAB_180006396	
goto LAB_18000643c;	CspBuildRootCertStore	18000d2ec	undefined CspBui		524		18000637c 4c 8b 0d	MOV	R9,qword ptr [WPP_pszIndent]	= NaP
3	CspBuildCertificateFilena.	18000dbb8	undefined CspBui		737		a5 6b 02 00			
if ((*(int *)(param_1 + 0x28) != 2)    (*(longlong *)(param_1 + 0x20) != 0)) {	CspBeginTransaction	18000bf7c	ulonglong CspBeg		279		180006383 8d 57 3a	LEA	param_2,[RDI + 0x3a]	
uVar2 = 0x57;	CspAuthenticateUser	180002a64	ulong CspAuthent		427		180006386 48 8b 49 10	MOV	param_1,qword ptr [param_1 + 0x10]=>DAT_	18002c = 000000030000000h
goto LAB_18000643c;	CspAuthenticatePin	18000ad70	undefined CspAut		726		18000638a 4c 8d 05	LEA	param_3,[WPP_e8978bfc35ec3505198246492c0	09beb = FCh
3	CspAppi yPKCS 15igningFor	180018818	undefined8 CspAp		652		2f dd 01 00			
ppcVar4 = *(char ***) (param_1 + 0x10);	CopAddCardCacheItem	180001560	void - CspAlloch		270		180006391 e8 2a b0	CALL	WPP_SF_s	undefined WPP_SF_s (ur
uVar3 = VerifyPKCS2Padding((longlong)ppcVar4,*(uint *)(param_1 + 0x18), <local_res20,local_res8);< td=""><td>CapAutocal ucachettem</td><td>100006000</td><td>arongrong cspAdd</td><td></td><td>2/0</td><td></td><td>ff ff</td><td></td><td></td><td></td></local_res20,local_res8);<>	CapAutocal ucachettem	100006000	arongrong cspAdd		2/0		ff ff			
uVar2 = (ulong)uVar3;										
ppcVar5 = local_res20;								LAB_180006396	XREF[3]	: 18000636e(j), 180006374
if (uVar2 == 0) {										18000637a(j)
uVar8 = local_res8[0] - 1;						→	180006396 48 85 db	TEST	RBX, RBX	-
if (uVar8 != 0) (							180006399 Of 84 8b	JZ	LAB_18000642a	
	1						00 00 00			_
console - scripting							18000639f 4d 85 ff	TEST	R15,R15	
							1800063a2 Of 84 82	JZ	LAB_18000642a	-
							00 00 00			
							1800063a8 4d 85 f6	TEST	R14,R14	

LAB\_18000642a dword ptr [RBX],0x2

LAB\_1800063bc

LAB\_18000643c

EBX, 0x51a

1800063ab 74 7d JZ

1800063ad 83 3b 02 CMP 1800063b0 74 0a

1800063b2 bb 1a 05

00 00 1800063b7 e9 80 00 JZ

MOV

JMP

### Windows Non Constant-time Protocol Code

<pre>1 int CPImportKey(csInfo) {</pre>						
2	•••					
3	<pre>int result = CspUnpadData(csInfo, &amp;len, &amp;out);</pre>					
4	<pre>if (result == SUCCESS) {</pre>					
5						
6	<pre>int isLegal = FIsLegalKeySize(len, out);</pre>					
7	<pre>if (isLegal) {</pre>					
8	<pre>Decrypt_3DES(encData, out, len, &amp;decData);</pre>					
9						
10	}					
11	}					
12	return SUCCESS;					
13 }						

# Our Target

- Nonconstant-time code in unpadding layer
- Flush+Reload cache attack
- Very noisy



### Bleichenbacher Attack – What Do We Need?

- Padding Oracle with feedback
- Ability to modify the ciphertext
- Large number of queries
- Large number of queries and FAST our deadline is 10 hours!

### How Do We Generate Messages?







**PIN** needed





# Harmless Looking Malicious HTML



Browsers can access remote files!

• Stealth!

• Repetition!

1	<html></html>
2	<head></head>
3	<meta content="4" http-equiv="refresh"/>
4	
5	<iframe src="file://win-r3f0hi0ntca\mysecret\&lt;/th&gt;&lt;/tr&gt;&lt;tr&gt;&lt;th&gt;&lt;/th&gt;&lt;th&gt;classifiedsecrets.txt"></iframe>
6	
7	

### Stealth Server Side

- Intercept and replace RSA message
- Server can detect the large volume of authentication requests
- Unauthenticated Data in Kerberos messages
- Response packet can be completely spoofed!
- We can create a fast and stealthy attack (no need for the KDC)







# Fast Enough?

- Short answer -> sometimes for some messages
- Real tokens (40k dataset generated by the KDC)
  - ~23% decrypted in less than 20K queries
  - ~8.5% decrypted in less than 10K queries



### Impersonation

- Decrypted tokens give the attacker access to the user/admin's resources
- We only need one token







# Bleichenbacher's Attack (1998)

- Padding structure constraint
  - x
- Initial multiplier search
  - X
- Range reduction Additional Multipliers



# Noisy Oracle

- Oracles are often noisy!
- FN and FP have different impacts
- Repeat all queries? Inefficient
- Attack can handle FN -> aim for lower FP probability
- Repetition of positives
  - Double repeat positive queries twice require both be positive
  - Majority repeat positive queries up to 3 times require majority be positive
- Can we improve this?



# Noisy Oracle

- Our oracle isn't binary!
  - Number of hits in our F+R attack matter!
- FP usually have just one hit
- Tailored query repetition algorithm



### Accelerated Attack- Experiment Results

- Focus on "Fast subset" 18% of messages
- Attack performed successfully on 15 of 17 "fast" tokens
- We still have a nonzero probability for a FP



### False Positive Detection

- Upper bound for amount of false queries
- Too many false queries in a rov
  -> FP
- State has changed
- Traceback!



### Traceback – Attack Recovery

- Save previous states
- Count false queries
- Revert to a previous state if FP is detected
- Traceback method much closer to the perfect oracle



### Who Is Fast?

- Correlation between attack steps
- Fast initial search -> Fast range reduction
- Message classification
- Detect early and abort!



### Early Abort – Experiment Results

- Experiment time: 3 days
- Detection limit: 600 queries
- Continue attack iff initial multiplier search is completed

#	$n_{msgs}$	$T_q$	T[hours]	$Q_{<16k}$
1	36	29336	18.4	8309
2	30	26185	16.4	8680
3	29	27023	16.9	9367

### Ghost of Crypto Past

- Nonconstant time implementations -> side channel attacks
- Lack of Forward Secrecy -> compromise other protocols
- Unlimited RSA decryption attempts why?



### Recap + Questions

- Broke Kerberos authentication protocol
- Improved Bleichenbacher Attack Methodology
- End-to-End impersonation attack on Windows
- One PIN code entry, one site, one day -> control!





Git Link